KeYmaera X

A Tutorial on Interactive Verification for Hybrid Systems

Nathan Fulton Marktoberdorf 2017

August 11, 2017



Examples: https://nfulton.org/marktoberdorf.zip Slides: https://nfulton.org/slides/marktoberdorf.pdf

Motivation

KeYmaera X provides strong evidence that Cyber-Physical Systems are safe. But you need to provide the model and sometimes help the proof.

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- Differential Dynamic Logics
 Syntax and Semantics
- Sound and relatively complete axiomatizations
- Some examples

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This Lecture:

- Practical advice for modeling systems
- Hands-on Exercise proving theorems
- Example-driven

Outline

Straight Line Dynamics The Stop Sign Model

Circular Dynamics

Loitering Outside Prohibited Airspace

Logarithmic Dynamics

Safe SCUBA Diving

Extras

The ODE Solver Taylor Approximations as Successive Differential Cuts The Stop Sign Model





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- KeYmaera X's edit tool checks your arithmetic (common and annoying source of errors, both in proofs and implementations!)
- Quantifier Elimination is a **powerful tool** useful for more than just decision procedures:
 - Find assumptions and loop invariants by reducing the system to arithmetic and eliminating quantifiers.
 - ► ModelPlex: ∀x₀, x₁,..., x_n.∃y₀,..., x_n.φ is kinda hard to check at runtime...

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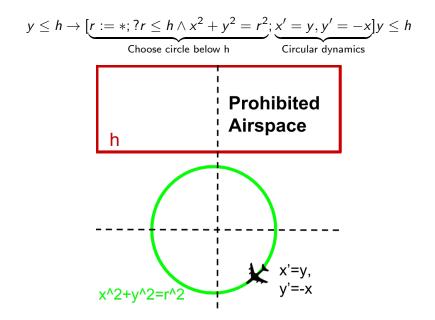
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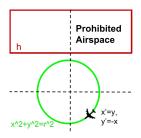
Loitering Outside Prohibited Airspace



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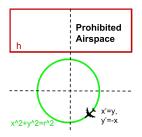


$$(y \le h)' \equiv (y)' \le (h)' \equiv -x \le 0$$
 FALSE



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$$(x^2+y^2=r^2)' \equiv (x^2+y^2)' = (r^2)' \equiv 2xx'+2yy' = 0 \equiv 2xy-2xy = 0$$



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$$r \leq h \wedge x^2 + y^2 = r^2 \rightarrow^? y \leq h$$

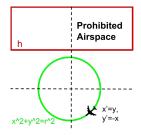


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 FALSE

COUNTER-EXAMPLE: $-2 \le -2 \land 3 + 1 = 4 \nrightarrow -1 \le -2$



On Annoying Assumptions

HOME > EXTREME > THE ESA HAS FIGURED OUT WHAT KILLED THE SCHIAPARELLI MARS LANDER

The ESA has figured out what killed the Schiaparelli Mars lander

By Jessica Hall on November 30, 2016 at 2:00 pm 38 Comments



The Schiaparelli lander model

When the navigation system got wind of the IMU's wacky output, it decided that meant the spacecraft had "an estimated altitude that was negative" — that is, below ground level. In its scramble, the system released the backshell too early, fired the braking thrusters, and finally flicked on the on-ground systems as if Schiaparelli had already

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- Like loop invariants, differential invariants sometimes need strengthening.
- In these cases, try using differential cuts to describe geometric constraints on the system.
- Most early proof attempts fail due to missing obvious assumptions:
 - Upper/lower-bounds (esp. positivity).
 - Missing t' = 1 in time-triggered systems.
 - Missing control epsilon $t \leq T$ in evolution domain.
 - Interesting dynamics (e.g., missing $v \ge 0$).

Use counter-examples to find these errors.

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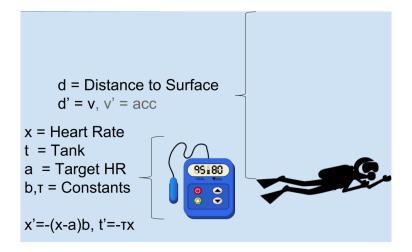
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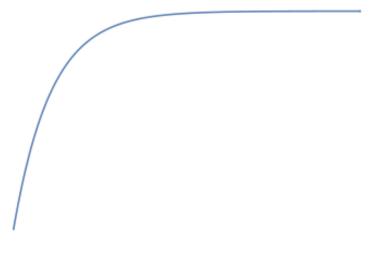
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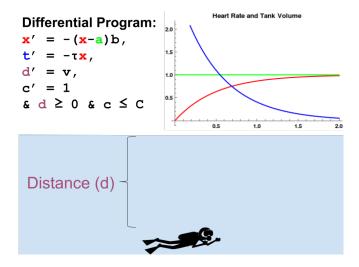


Heart Rate Function



 $x' = -(x - HR_{max})b$

SCUBA Ascent Case



Control Goal: Find a condition that ensures the diver reaches the surface before running out of oxygen.

$$x' = -(x - a)b, t' = - au x, d' = v, c' = C$$
 & $c \le C \land d \ge 0$

Idea: Bound time and all non-linear terms, then prove linear inequalities on these bounds by integrating.

$$x'=-(x-a)b,t'=- au x,d'=v,c'=C$$
 & $c\leq C\wedge d\geq 0$

Idea: Bound time and all non-linear terms, then prove linear inequalities on these bounds by integrating.

▶ Non-linear term: $x \leq HR_{max}$

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- ▶ Bound time: $d_0 + vc \ge 0 \Rightarrow$ bound on time (denote as $z = \frac{-d}{v_0}$).

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$$t = t_0 - \tau xc \ge t_0 - \tau HR_{max}c \ge \underbrace{t_0 - \tau HR_{max}z \ge 0}_{\text{lottal safe states!}}$$

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$$t = t_0 - \tau xc \ge t_0 - \tau HR_{max}c \ge \underbrace{t_0 - \tau HR_{max}z \ge 0}_{\text{Initial safe states!}}$$

The first step requires $x \leq HR_{max}$. This is the only interesting lemma.

Let's prove $x < HR_{max}$ instead to avoid extra case splitting due to the $x = HR_{max}$ bifurcation point.

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Step 1: Find an existential condition equivalent to our goal:

 $\models_{\mathbb{R}C\mathcal{F}} x < HR_{max} \leftrightarrow \exists y.?$

Let's prove $x < HR_{max}$ instead to avoid extra case splitting due to the $x = HR_{max}$ bifurcation point.

Step 1: Find an existential condition equivalent to our goal:

$$\models_{\mathbb{R}C\mathcal{F}} x < HR_{max} \leftrightarrow \exists y. \ y^2(x - HRmax) = -1$$

Let's prove $x < HR_{max}$ instead to avoid extra case splitting due to the $x = HR_{max}$ bifurcation point.

Step 1: Find an existential condition equivalent to our goal:

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Step 2: Find y' s.t. $(y^2(x - HR_{max}) = -1)'$ is true:

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► Step 2: Find y' s.t.
$$(y^2(x - HR_{max}) = -1)'$$
 is true:
 $(y^2(x - HRmax) = -1)' \equiv (y^2(x - HR_{max}))' = 0$
 $\equiv 2yy'(x - HR_{max}) + y^2x' = 0$
 $\equiv 2yy'(x - HR_{max} + y^2(-(x - a)b) = 0$
 $\equiv \dots$
 $\equiv y' = \frac{b}{2}y$

(All equivalences are with respect to the ODE.)

Take-aways from SCUBA Example

- As systems become harder to model, parametric models save the day.
- Identifying and using differential ghosts is (sometimes) systematic.
- Partial solutions to fragments of an ODE's dynamics are useful whenever you can upper-bound terms.
- ► Tactics ⇒ proof reuse

Summary

STOP







Resources

Notes, slides, and examples from this talk:

https://nfulton.org/marktoberdorf

KeYmaera X website:

https://keymaeraX.org

Online Instance (With Mathematica!):

https://web.keymaeraX.org

Source Code (Scala):

https://github.com/LS-Lab/KeYmaeraX-release

KeYmaera X Credits: Stefan Mitsch, Jan-David Quesel, Marcus Völp, Brandon Bohrer, Yong Kiam Tan, André Platzer, ... **SCUBA Credits:** Karim Elmaaroufi and Viren Bajaj

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Add a time variable:

$$[x' = v, v' = a, t' = 1]P(x, v)$$

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Use differential cuts to add solutions in linear order:

$$[x' = v, v' = a, t' = 1 \& v = at + v_0 \land x = \frac{at^2}{2} + v_0 t + x_0]P(x, v)$$

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Rewrite the post-condition in terms of t:

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Inverse differential ghosts to remove all equations except time:

$$[t' = 1\&v = at + v_0 \land x = \frac{at^2}{2} + v_0t + x_0]P(t)$$

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Use univariate solve:

$$orall sorall 0 \leq t \leq s.v = at + v_0 \wedge x = rac{at^2}{2} + v_0t + x_0
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Taylor Approximations in KeYmaera X

$$s'=c, c'=-s$$

$$s = \sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \dots$$

$$c = \cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \dots$$

Taylor Approximations in KeYmaera X

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