

**15-819M Data, Code, Decisions**  
**Assignment 1**      ( $\Sigma 50$ )      **due by Tue 9/29/2009**

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Disclaimer: No solution will be accepted that comes without an **explanation!**

**Exercise 1      Propositional Sequent Calculus (28p)**

1. Which of the following formulas is true? Give a counterexample or a proof in sequent calculus.
  - a)  $(a \leftrightarrow b) \leftrightarrow (a \rightarrow \neg b) \wedge (\neg b \rightarrow a)$
  - b)  $(c \leftrightarrow d) \rightarrow ((a \leftrightarrow c) \leftrightarrow (a \leftrightarrow d))$
  - c)  $\neg a \rightarrow (a \rightarrow b)$
  - d)  $\left( \left( (p \rightarrow q) \rightarrow (\neg r \rightarrow \neg s) \right) \rightarrow r \right) \rightarrow \left( (t \rightarrow p) \rightarrow (s \rightarrow p) \right)$
2. In at least two of the above cases, please send a KeY<sup>1</sup> proof file by email.
3. Give sequent proof rules for the following operators:
  - a) XOR (exclusive-or)
  - b) NAND (negated and)
  - c) NOR (negated or)
  - d)  $A?B : C$  with the semantics

$$\llbracket A?B : C \rrbracket_I = \begin{cases} \llbracket B \rrbracket_I & \text{if } \llbracket A \rrbracket_I = \text{true} \\ \llbracket C \rrbracket_I & \text{if } \llbracket A \rrbracket_I = \text{false} \end{cases}$$

**Exercise 2      Propositional Proofs (10p)**

Is the following rule a replacement for the implication left rule?

$$\frac{\Gamma \vdash A, \Delta \quad \Gamma, A, B \vdash \Delta}{\Gamma, A \rightarrow B \vdash \Delta}$$

What is the advantage of this rule? Is it a sound replacement? Is it a complete replacement for the implication left rule?

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<sup>1</sup> <http://www.key-project.org/download/releases/webstart/KeY.jnlp>

### Exercise 3 DPLL (12p)

1. Prove or disprove unsatisfiability of the conjunction of the following formulas with DPLL

$$\neg a \vee \neg b \vee c \vee \neg d$$

$$\neg b \vee \neg c$$

$$b \vee \neg c \vee d$$

$$\neg a \vee c$$

$$\neg a \vee b \vee \neg d$$

$$a \vee d$$

2. If the previous formulas are unsatisfiable, what is the simple-most formula to add or to remove to make the formulas satisfiable? If the previous formulas are satisfiable, what is the simple-most formula to add or to remove to make the formulas unsatisfiable? In each case, prove or disprove in DPLL.